A locally finite tree that behaves like an infinite star

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Abstract

We construct a tree T of maximal degree 3 with infinitely many leaves such that whenever finitely many of them are removed, the remaining tree is isomorphic to T. In this sense T resembles an infinite star.

1 Introduction and general remarks

The famous Rado graph R (see [2],[3]) has the property of being isomorphic to R-W, where $W\subset V(R)$ is any finite set of vertices. A similar property of a countable tree T with infinitely many leaves would be $T\cong T-W$ for any finite set of leaves W. Trivially, this property is satisfied by an infinite star, but what happens if we restrict ourselves to locally finite trees? If such a tree exists, it must be fairly "large". In this article we construct such a tree T of maximal degree 3, which is best possible.

The problem is motivated by the work of Bonato and Tardif [1] and the author [4] on mutual embeddings of infinite trees. Call two trees twins or mutually embeddable if each of them contains an isomorphic copy of the other as a subtree. Define the twin number of a tree T, in notation m(T), to be the number of isomorphism classes of trees twinned with T. Bonato and Tardif [1] asked, what values the twin number can take, i.e. how many pairwise non-isomorphic trees can be mutually embeddable. Their conjecture was that the twin number of a tree is always either 1 or infinite. In [4] the author considered this problem for the class of locally finite trees. While it is easy to construct locally finite trees of twin number 1, we could not find a locally finite tree T of m(T) = 1 which has a non-surjective self-embedding, other than the one-way infinite path. It can be shown that such a tree must be isomorphic to T-x for infinitely many of its leaves x, so the natural question to ask is, if such a tree exists. The aim of this paper is to give an affirmative answer, and in fact the statement we prove is much stronger. Our construction might be helpful for solving the original problem and, we believe, it should be interesting in its own right.

The additional restriction of having infinitely many leaves is not substantial, as the only tree with finitely many leaves and the above property is a ray, i.e. a one-way infinite path. More dramatically, a ray is the only tree T with finitely many leaves satisfying $T \cong T - x$ for *some* leaf x. Indeed, the removal of a leaf

from a tree with finitely many leaves either decreases the number of leaves by 1 or does not affect it. In the latter case the vertex next to the removed leaf must have had degree 2. Therefore, if T has a vertex of degree at least 3 and we start removing leaves of T one by one in any order, we will at some point obtain a tree with fewer leaves than originally. In particular, repeated application of the isomorphism between T and T-x would result in a tree with fewer leaves than T but isomorphic to T, a contradiction.

We now describe the construction of the desired tree T. First of all, observe that it suffices to ensure $T \cong T - x$ for any single leaf x, for then the desired property follows by iteration. We build up T inductively as a union of an ascending chain of trees $T_0 \subset T_1 \subset T_2 \subset \cdots$ with $\Delta(T_n) = 3$ for all n. To be precise, we set T = (V, E) where $V = \bigcup_{i=0}^{\infty} V(T_i)$, noting that the sequence $\{V(T_i)\}_{i=0}^{\infty}$ is ascending as well, and let $e \in E$ if and only if $e \in E(T_i)$ for some i. It is immediate that T is another tree of maximal degree 3.

Let the core c(T) of a tree T be the graph spanned by all vertices v such that T-v has at least two infinite components. It is not hard to see that the core must be connected, i.e. c(T) is a subtree of T. Note also, that the core is invariant under any isomorphism $\phi \colon T \to T-x$, i.e. ϕ restricted to c(T) defines an automorphism of c(T), which we also denote by ϕ . Consequently, the inverse map $\phi^{-1} \colon T-x \to T$ gives rise to the inverse automorphism $\phi^{-1} \colon c(T) \to c(T)$.

Let S and T be two trees with vertices $v \in S$ and $w \in T$. The sum of two rooted trees (S, v) + (T, w) is constructed by identifying v with w and "gluing together" S and T accordingly. Similarly, if $W \subset V(T)$, we can define (T, W) + (S, v) by attaching a copy of S to each $w \in W$ identifying v with w.

Let x_1, x_2, \ldots, x_n be leaves of a tree T such that $T \cong T - x_i$ for each i and let ϕ_i be the corresponding isomorphism between T and $T - x_i$. Denote by Φ the set $\{\phi_1, \phi_2, \ldots, \phi_n\}$. For $Y \subset c(T)$, define the closure $\Phi(Y)$ of Y under Φ as the set of all vertices $w \in c(T)$ that can be mapped to some $y \in Y$ by finitely applications of ϕ_1, \ldots, ϕ_n and $\phi_1^{-1}, \ldots, \phi_n^{-1}$, that is, $\Phi(Y)$ is the image of Y under $\langle \Phi \rangle$, the free group generated by Φ . Furthermore, given a rooted tree (S, s) we define $*(T, Y, S, s; \Phi)$, the convolution of (T, Y) and (S, s) via Φ , to be the tree $(T, \Phi(Y)) + (S, s)$. As $Y \subset \Phi(Y)$, the convolution $*(T, Y, S, s; \Phi)$ extends the sum (T, Y) + (S, s). More importantly, $*(T, Y, S, s; \Phi) - x_i \cong *(T, Y, S, s; \Phi)$ for all i such that the underlying isomorphism extends the one between T and $T - x_i$ and conversely, the convolution $*(T, Y, S, s; \Phi)$ is the minimal extension of (T, Y) + (S, s) preserving the above isomorphism.

The construction of the desired tree T comprises two major parts. First we construct the sequence $\{T_n\}$. Then we show that its union T has the needed property.

2 Constructing the sequence

Our aim is to construct an ascending sequence of trees $T_0 \subset T_1 \subset T_2 \subset \cdots$ along with a sequence of vertices x_0, x_1, x_2, \ldots such that for all n the following conditions are satisfied:

1.
$$\Delta(T_n) = 3$$
.

- 2. x_i is a leaf of T_n for all $i \leq n+1$.
- 3. $T_n \cong T_n x_i$ for all $i \leq n$. Moreover, for $i \leq m < n$ the isomorphism between T_n and $T_n x_i$ extends the isomorphism between T_m and $T_m x_i$.
- 4. Writing ϕ_i^n for the isomorphism between T_n and $T_n x_i$, as in 3. and Φ for the set $\{\phi_0^n, \ldots, \phi_n^n\}$, there exists a vertex $y \in c(T_n)$ such that each vertex $y' \in \Phi(y)$ has degree 2 in T_n .
- 5. Among all leaves in T_n apart from x_0, x_1, \ldots, x_n the leaf x_{n+1} has the minimal distance from x_0 .

Let T_0 be a doubly infinite ray, labeled by the integers ..., -1, 0, 1, 2, ... with a leaf attached to each *even* non-negative vertex. If we let x_0 and x_1 be the leaves attached to the vertices labeled by 0 and 2, then a trivial check confirms properties 1–5 for T_0 .

Suppose that we have constructed T_n along with x_{n+1} . We wish to extend T_n to T_{n+1} . Note that, by property 3 and the fact that T_n is not a ray (which is implied by property 1) T_n must have infinitely many leaves, as explained in the first section. Thus, once we have constructed T_{n+1} such that properties 1–4 are satisfied, we can choose x_{n+2} according to 5.

We want to construct T_{n+1} as a union of another ascending sequence of trees $U_0 \subset U_1 \subset U_2 \subset \cdots$ where $U_0 = T_n$ and for all i the tree U_i should have the following properties:

- (i) $\Delta(U_i) = 3$.
- (ii) $x_0, x_1, \ldots, x_{n+1}$ are leaves of U_i .
- (iii) If i is even, then $U_i \cong U_i x_j$ for j = 0, 1, ..., n. If i is odd, then $U_i \cong U_i x_{n+1}$. In both cases the corresponding isomorphism between U_i and $U_i x_j$ extends the one between U_{i-2} and $U_{i-2} x_j$.

At the same time we construct a sequence of vertex sets $W_i \subset U_{i-1}$, whose meaning will become clear later on.

Choose a vertex y as in 4. Attach a new path P of length 2 to y, whose other end we denote by y'. Attach a new doubly infinite path R to y'. Let R_1 and R_2 denote the two rays, into which y' divides R. Now to each vertex on R_1 except y' attach a copy of $(P, y) + (T_n, y)$ in the same way as it is attached to y'. To each vertex on R_2 attach a copy of $(P, y) + (T_n - x_{n+1}, y)$. Call the resulting tree U_1 and let $W_1 = \{y\}$ and $S = U_1 - (U_0 - y)$. In other words, $U_1 = (U_0, y) + (S, y)$.

To obtain U_{2i} from U_{2i-1} , set $U_{2i} = *(U_{2i-2}, W_{2i-1}, S, y; \Phi_{2i-2})$ where Φ_{2i-2} is the set of isomorphisms removing x_0, \ldots, x_n from U_{2i-2} (note that $U_{2i-1} = (U_{2i-2}, W_{2i-1}) + (S, y)$ so U_{2i} extends U_{2i-1}). Define W_{2i} to be $\Phi_{2i-2}(W_{2i-1}) \setminus W_{2i-1}$, in other words W_{2i} is the set of those vertices in U_{2i-1} to which we attached a new copy of (S, y).

Similarly, to obtain U_{2i+1} from U_{2i} , set $U_{2i+1} = *(U_{2i-1}, W_{2i}, S, y; \Phi_{2i-1})$ where Φ_{2i-1} consists of the single isomorphism between U_{2i-1} and $U_{2i-1} - x_{n+1}$. Again, $U_{2i} = (U_{2i-1}, W_{2i}) + (S, y)$, and hence U_{2i+1} extends U_{2i} . Define W_{2i+1} to be $\Phi_{2i-1}(W_{2i}) \setminus W_{2i}$.

Note that $W_1 \subset c(U_0)$ and therefore, by the invariance of the core, $W_2 \subset \Phi_0(W_1) \subset c(U_0) \subset c(U_1)$. Hence, $W_3 \subset \Phi_1(W_2) \subset c(U_1) \subset c(U_2)$. It follows by induction that $W_i \subset c(U_{i-1})$ for all i. This shows that a leaf of U_i remains a leaf in U_{i+1} . In particular, $x_0, x_1, \ldots, x_{n+1}$ are leaves of all U_i and (ii) holds.

Another important observation is the fact that, for even i, the set W_i lies on the same side of P as y, whereas for odd i > 3, W_i lies on the same side as y'.

We also note that property (iii) is immediate from the construction and the corresponding property of the convolution.

We now must prove (i). For U_1 this follows by construction. Property 4 of T_n gives us $\Delta(U_2)=3$. To prove the general statement it suffices to show that $d_{U_i}(w)=2$ for all $w\in W_{i+1}$. Indeed, since we know that $d_S(y)=1$, this will imply that at each step we identify vertices of degree 2 with vertices of degree 1, i.e. no vertex of degree 4 or greater is generated. If $w\in W_{i+1}$, then $w=\phi(w')$ for some $w'\in W_i$ and $\phi\in \langle\Phi\rangle$ (group generated by Φ). Therefore, by the induction hypothesis, $d_{U_i}(w)=d_{U_i}(w')=2$ unless $w\in \Phi_{i-1}(v)$ for some v adjacent to one of x_0,\ldots,x_n if i is odd and adjacent to x_{n+1} if i is even. Note, however, that $W_3\not\subset U_0$, which means that for odd i we have $W_i\not\subset U_0$, so w or w' cannot be mapped to a vertex adjacent to x_0,\ldots,x_n . And if i is even, the above obstruction can only occur in the case i=2, however, by construction of U_1 we know that this cannot happen. Therefore, property (i) holds.

We need to show that T_{n+1} , defined as the union of $U_0 \subset U_1 \subset \cdots$, satisfies 1-4. Property 1 is immediate from (i). Property 2 follows from (ii) and the fact that we can always choose x_{n+2} according to 5. Property 3 is a consequence of (iii). So the only property left to be checked is 4.

Let z be the center of P, i.e. the vertex between y and y'. Note that, by construction of T_{n+1} , vertices in $\Phi(z)$ cannot have degree 1 in T_{n+1} . If some $z' \in \Phi(z)$ has degree 3 then z' has degree 3 in some U_i . But this implies that z has degree 3 in some U_j , which contradicts the fact that $z \notin W_j$ for any j.

3 Taking the union

We now define T to be the union of the ascending sequence $\{T_n\}$. We already mentioned that $\Delta(T)=3$. Note also that, by property 3 of the sequence $\{T_n\}$, x_i is a leaf of T for all i and property 5 implies that T has no other leaves. Finally, by property 3, T is isomorphic to $T-x_i$ for all i. Indeed, since the isomorphisms $\phi_i^n: T_n \to T_n - x_i$ are "nested", we can combine them and define $\phi_i: T \to T - x_i$ by $\phi_i(y) = \phi_i^n(y)$, where n is the smallest index satisfying $y \in T_n$. Hence, we have shown that T is a tree with the desired properties.

4 Further remarks

The above construction seems to have a very high "degree of freedom", i.e. altering the construction one can obtain many pairwise nonisomorphic trees with the above property. It would be interesting to find out what properties a locally finite tree can have in addition to T with $T\cong T-x$ for each leave x. In particular, we do not know, whether such a tree T must or can be isomorphic to T-X for some

infinite set of leaves X. This is closely related to the problem of twin numbers and the solution could shed more light on "paradoxical" properties of graphs.

References

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